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Three Datatype Defining Rewrite Systems for Datatypes of Integers each extending a Datatype of Naturals (version 3)

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Three Datatype Defining Rewrite Systems for Datatypes of Integers each extending a Datatype of Naturals*

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Abstract

Integer arithmetic is specified according to three views: unary, binary, and decimal notation. In each case we find a ground-confluent and terminating datatype defining rewrite system. In each case the resulting datatype is a canonical term algebra which extends a corresponding canonical term algebra for natural numbers. For each view, we also consider an alternative rewriting system.

Keywords and phrases: Equational specification, initial algebra, datatype defining rewrite system, abstract datatype.

Contents

*Version 3: Some new non-confluence and termination results recorded in [12] are mentioned and some errors recorded in [12] are corrected; ground-completeness of $\mathbb{Z}_{bud}$ (Tables 7 and 8, pages 7, 8) is proven; a DDRS for the ring of Integers is added (Table 18, pages 18, 19) and its ground-completeness is proven.

$0' = 1$	$3' = 4$	$6' = 7$
$1' = 2$	$4' = 5$	$7' = 8$
$2' = 3$	$5' = 6$	$8' = 9$

Table 1: Enumeration and successor notation of digits of type $\mathbb{Z}$

1 Introduction

Using the specifications for natural numbers from [1] we develop specifications for datatypes of integers. We will entertain the strategy of [1] to develop different views characteristic for unary notation, binary notation, and decimal notation respectively. Each of the specifications is a so-called DDRS (datatype defining rewrite system) and consists of a number of equations that define a term rewriting system by orienting the equations from left-to-right. A DDRS, or more precisely, the associated term rewriting system, must be strongly terminating and ground-confluent.

This paper is a sequel to the report [1] which deals with DDRSes for the natural numbers and it constitutes a further stage in the development of a family of arithmetical datatypes with corresponding specifications. The resulting specifications (DDRSes) incorporate different "views" on the same abstract datatype. The *unary view* provides a term rewriting system where terms in unary notation serve as normal forms. The unary view also provides a semantic specification of binary notation, of decimal notation, and of hexadecimal notation. The three logarithmic notations were modified in [1] with respect to conventional notations in such a way that syntactic confusion between these notations cannot arise. In this paper, the *hexadecimal view* is left out as that seems to be an unusual viewpoint for integer arithmetic.

It seems to be the case that for the unary view the specification of the integers (given in Table 3) is entirely adequate, whereas all subsequent specifications for binary view and decimal view may provide no more than a formalization of a topic which must be somehow understood before taking notice of that same formalization. It remains to be seen to what extent the first DDRS for the unary case may serve exactly that expository purpose.

The strategy of the work is somewhat complicated: on the one hand we look for specifications that may genuinely be considered introductory, that is, descriptions that can be used to construct the datatype at hand for the first time in the mind of a person. On the other hand awareness of the datatype in focus may be needed to produce an assessment of the degree of success achieved in the direction of the first objective.

1.1 Digits and rewrite rules in equational form

Digits are $0, 1, 2, 3, 4, 5, 6, 7, 8, 9$, and are ordered in the common way:

$$0 < 1 < 2 < 3 < 4 < 5 < 6 < 7 < 8 < 9.$$

For the digits $0, 1, 2, 3, 4, 5, 6, 7, 8$ we denote with i' the successor digit of i in the given enumeration. In Table 1 the successor notation on digits is specified as a transformation of syntax, and we adopt this notation throughout the paper.

We will list rewrite rules in the form of equations $t = r$ to be interpreted from left-to-right,

and we will add tags of the form

$$[\mathrm{Nn}] \quad t = r$$

for reference, with "N" some name and "n" a natural number (in ordinary, decimal notation). Furthermore, for $k, \ell \in \{0,1,2,3,4,5,6,7,8,9\}$ and $k < \ell$, the notation

$$[\mathrm{Nn}.i]_{i=k}^{\ell} \quad t = r$$

represents the following $\ell - k + 1$ equations:

$$[\mathrm{Nn}.k] \quad t[k/i] = r[k/i], \quad \ldots \quad , [\mathrm{Nn}.\ell] \quad t[\ell/i] = r[\ell/i],$$

thus with i instantiated from k to ℓ. Occasionally, we will use this notation with two "digit counters", as in

$$[\mathrm{Nn}.i.j\,]_{i,j=k}^{\ell} \quad t = r,$$

for a concise representation of the following $(\ell - k + 1)^2$ equations:

$$[\mathrm{Nn}.k.k] \quad t[k/i][k/j] = r[k/i][k/j], \quad \ldots \quad , [\mathrm{Nn}.\ell.\ell] \quad t[\ell/i][\ell/j] = r[\ell/i][\ell/j].$$

1.2 A signature for integers

The signature $\Sigma_{\mathbb{Z}}$ has the following elements:

1. A sort $\mathbb{Z}$,
2. For digits the ten constants $0, 1, 2, 3, 4, 5, 6, 7, 8, 9$,
3. Three one-place functions $S, P, - : \mathbb{Z} \to \mathbb{Z}$, "successor", "predecessor", and "minus", respectively,
4. Addition and multiplication (infix) $+, \cdot : \mathbb{Z} \times \mathbb{Z} \to \mathbb{Z}$,
5. Two one-place functions (postfix) $\_ :_b 0, \_ :_b 1 : \mathbb{Z} \to \mathbb{Z}$, "binary append zero" and "binary append one", these functions will be used for binary notation,
6. Ten one-place functions (postfix)

   $$\_ :_d 0, \_ :_d 1, \_ :_d 2, \_ :_d 3, \_ :_d 4, \_ :_d 5, \_ :_d 6, \_ :_d 7, \_ :_d 8, \_ :_d 9 : \mathbb{Z} \to \mathbb{Z},$$

   "decimal append zero", ...,"decimal append nine", to be used for decimal notation.

The "append <digit name>" functions defined in items 5 and 6 can be viewed as instantiations of more general two-place "append" functions, but that would require the introduction of sorts for bits (binary digits) and for decimal digits. However, we prefer to keep the signature single-sorted and that is why we instantiate such "digit append" functions per digit to unary functions and why we use postfix notation for applications of these functions. E.g.,

$$(9 :_d 7) :_d 5 \quad \text{and} \quad ((1 :_b 0) :_b 0) :_b 1$$

represent the decimal number 975, and the binary number 1001, respectively.

For the unary view the normal forms are the classical successor terms, that is

$$0, S(0), S(S(0)), ...,$$

[S1]	$x + 0 = x$	$[\mathrm{S5}.i]_{i=0}^{8}$	$i' = S(i)$
[S2]	$x + S(y) = S(x + y)$	$[\mathrm{S6}.i]_{i=0}^{1}$	$x :_b i = (x \cdot S(1)) + i$
[S3]	$x \cdot 0 = 0$	$[\mathrm{S7}.i]_{i=0}^{9}$	$x :_d i = (x \cdot S(9)) + i$
[S4]	$x \cdot S(y) = (x \cdot y) + x$		

Table 2: A DDRS for $\mathbb{N}_{ubd}$, natural numbers in unary view

and all minus instances $-(t)$ for each such nonzero normal form t, e.g. $-(S(S(0)))$. If no confusion can arise, we abbreviate $-(t)$ to $-t$, as in $-x$. As an alternative (and introduction to subsequent DDRSes) we shall briefly consider a DDRS that is based on a "unary append zero" function.

For the binary view and for the decimal view, we provide one DDRS for each. Normal forms are all appropriate digits, all applications of the respective append functions to a nonzero normal form, and all minus instances $-t$ for each such normal form t that differs from 0. Thus $-(((1 :_b 0) :_b 0) :_b 1)$ is an example of a normal form in binary view, and $-((9 :_d 7) :_d 5)$ is one in decimal view.

2 One ADT, three datatypes

An abstract datatype (ADT) may be understood as the isomorphism class of its instantiations which are datatypes. The datatypes considered in [1] are so-called canonical term algebras which means that carriers are non-empty sets of closed terms which are closed under taking subterms.

2.1 Unary view

Table 2 provides a DDRS for the natural numbers and defines the datatype $\mathbb{N}_{ubd}$. Minus and predecessor are absent in this datatype. Successor terms, that is expressions involving zero and successor only, serve as normal forms for the datatype $\mathbb{N}_{ubd}$. This DDRS contains the well-known equations [S1] – [S4] and the twenty-one equations $[\mathrm{S5}.i]_{i=0}^{8}$ – $[\mathrm{S7}.i]_{i=0}^{9}$, and defines the rewrite rules that serve the rewriting of binary and decimal notation.

In Table 3 a DDRS is provided of the integer numbers $\mathbb{Z}_{ubd}$ with successor, predecessor, addition, and multiplication, which are defined by equations [u1] – [u14]. We notice that we do not need equations for rewriting

$$(-x) \cdot y$$

because multiplication is defined by recursion on its right-argument, and that is why equation [u14] is sufficient, and why addition is defined by recursion on both its arguments and also requires [u11]. Like before, the twenty-one equations $[\mathrm{u15}.i]_{i=0}^{8}$ – $[\mathrm{u17}.i]_{i=0}^{9}$ serve the rewriting of binary and decimal notation.

In Table 4 one finds a listing of equations that are true in the datatype $\mathbb{Z}_{ubd}$ that is specified by the DDRS of Table 3. This ensures that these equations are semantic consequences of the equations for commutative rings.

So, binary and decimal notation are defined by expanding terms into successor terms. This expansion involves a combinatorial explosion in size and renders the specification in Tables 2 and 3 irrelevant as term rewriting systems from which an efficient implementation can be generated.

[u1]	$-0 = 0$	[u12]	$x \cdot 0 = 0$
[u2]	$-(-x) = x$	[u13]	$x \cdot S(y) = (x \cdot y) + x$
[u3]	$P(0) = -S(0)$	[u14]	$x \cdot (-y) = -(x \cdot y)$
[u4]	$P(S(x)) = x$	$[\text{u15}.i]_{i=0}^{8}$	$i' = S(i)$
[u5]	$P(-x) = -S(x)$	$[\text{u16}.i\,]_{i=0}^{1}$	$x :_b i = (x \cdot S(1)) + i$
[u6]	$S(-(S(x))) = -x$	$[\text{u17}.i]_{i=0}^{9}$	$x :_d i = (x \cdot S(9)) + i$
[u7]	$x + 0 = x$		
[u8]	$0 + x = 0$		
[u9]	$x + S(y) = S(x + y)$		
[u10]	$S(x) + y = S(x + y)$		
[u11]	$(-x) + (-y) = -(x + y)$		

Table 3: A DDRS for $\mathbb{Z}_{ubd}$, integer numbers in unary view

$$x + (y + z) = (x + y) + z \tag{1}$$
$$x + y = y + x \tag{2}$$
$$x + 0 = x \tag{3}$$
$$x + (-x) = 0 \tag{4}$$
$$(x \cdot y) \cdot z = x \cdot (y \cdot z) \tag{5}$$
$$x \cdot y = y \cdot x \tag{6}$$
$$1 \cdot x = x \tag{7}$$
$$x \cdot (y + z) = (x \cdot y) + (x \cdot z) \tag{8}$$
$$S(x) = x + 1 \tag{9}$$
$$P(x) = x + (-1) \tag{10}$$
$$x :_b i = (x + x) + i \qquad \text{for } i \in \{0, 1\} \tag{11}$$
$$x :_d i = (\underline{10} \cdot x) + \underline{i} \qquad \text{for } i \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}, \quad \underline{0} = 0,\ \ \underline{i'} = \underline{i} + 1,\ \ \underline{10} = \underline{9} + 1 \tag{12}$$

Table 4: Equations valid in $\mathbb{Z}_{ubd}$, where (1) – (8) axiomatize commutative rings

[u′1]	$x + 0 = x$	[u′3]	$x \cdot 0 = 0$
[u′2]	$x + (y :_u 0) = (x :_u 0) + y$	[u′4]	$x \cdot (y :_u 0) = (x \cdot y) + x$

Table 5: A DDRS for $\mathbb{N}_{u'}$, natural numbers in unary view with zero append

In the sequel we will consider specifications where normal forms are in binary notation and in decimal notation, respectively, that also employ the successor and predecessor functions. These specifications are far more lengthy and involved, but as DDRSes their quality improves because normal forms are smaller and are reached in fewer rewriting steps.

In order to give a smooth introduction to the subsequent DDRSes (for binary view and decimal view), we end with a brief exposition of a very simple alternative to the above DDRSes. Consider the extension of the signature $\Sigma_{\mathbb{Z}}$ defined in Section 1.2 with a one-place function (postfix)

$$\_ :_u : \mathbb{Z} \to \mathbb{Z},$$

the "unary append zero" function, or briefly: *zero append*, which can be used as an alternative for unary notation.

Normal forms in this alternative unary view are 0 for zero, and applications of the zero append function that define all successor values (each natural number n is represented by n applications of the zero append and can be seen as representing a sequence of 0's of length $n+1$). Furthermore, all minus instances $-t$ for each such normal form t that differs from 0 define the negative integers, so

$$-((0 :_u 0) :_u 0)$$

is an example of a normal form in this unary view (and it represents -2 in decimal notation).

Addition and multiplication are easy to define: Table 5 provides a DDRS for the natural numbers $\mathbb{N}_{u'}$. Termination follows with the use of a weight function, and also ground-confluence follows easily.

The DDRS that defines the extension of $\mathbb{N}_{u'}$ to integer numbers $\mathbb{Z}_{u'}$ is given in Table 6 below. It is immediately clear how rules for rewriting to unary notation with successor and predecessor, and binary and decimal notation can be defined, but we refrain from doing so and stick to the signature $\Sigma_{\mathbb{Z}}$ defined in Section 1.2 because the main purpose of the specifications in Table 5 and Table 5+6 is to introduce working with the "append functions", as we will do in our definitions of the datatypes in the subsequent sections.

[u′5]	$-0 = 0$	[u′9]	$(-(y :_u 0)) + (x :_u 0) = x + (-y)$
[u′6]	$-(-x) = x$	[u′10]	$(-x) + (-y) = -(x + y)$
[u′7]	$0 + x = x$	[u′11]	$x \cdot (-y) = -(x \cdot y)$
[u′8]	$(x :_u 0) + (-(y :_u 0)) = x + (-y)$	[u′12]	$(-x) \cdot y = -(x \cdot y)$

Table 6: Combined with Table 5, a DDRS for $\mathbb{Z}_{u'}$ that specifies integer numbers in unary view with zero append

$[\text{b1}.i]_{i=0}^{1}$	$0:_b i = i$	[b11]	$x \cdot 0 = 0$
[b2]	$S(0) = 1$	[b12]	$x \cdot 1 = x$
[b3]	$S(1) = 1:_b 0$	$[\text{b13}.i]_{i=0}^{1}$	$x \cdot (y:_b i) = ((x \cdot y):_b 0) + (x \cdot i)$
[b4]	$S(x:_b 0) = x:_b 1$	$[\text{b14}.i]_{i=1}^{8}$	$i' = S(i)$
[b5]	$S(x:_b 1) = S(x):_b 0$	$[\text{b15}.i]_{i=0}^{9}$	$x:_d i = (x \cdot S(9)) + i$
[b6]	$x + 0 = x$		
[b7]	$0 + x = x$		
[b8]	$x + 1 = S(x)$		
[b9]	$1 + x = S(x)$		
$[\text{b10}.i.j]_{i,j=0}^{1}$	$(x:_b i) + (y:_b j) = S^j((x + y):_b i)$		

Table 7: A DDRS for $\mathbb{N}_{bud}$, natural numbers in binary view

2.2 Binary view

In Table 7 a DDRS for a binary view of natural numbers is displayed, which employs the successor function as an auxiliary function. Leading zeros except for the zero itself are removed by $[\text{b1}.i]_{i=0}^{1}$, and successor terms are rewritten according to [b2] – [b5]. This DDRS contains fifteen (parametric) equations (that is, eighteen equations for the specification of addition and multiplication, and eighteen that serve the rewriting from decimal notation to binary notation via successor terms[1]). Furhermore, in equation $[\text{b10}.i.j]_{i,j=0}^{1}$ we use the abbreviation S^j for j applications of the successor function S, thus $S^0(t) = t$ and $S^{j+1}(t) = S(S^j(t))$. In the binary view natural numbers are identified with normal forms in binary notation. The specification has a canonical term algebra $\mathbb{N}_{bud}$ which is isomorphic to the canonical term algebra $\mathbb{N}_{ubd}$ of the specification in Table 2. In [12] it is shown that the rewriting system defined by this DDRS is complete.

In Table 8 minus and predecessor are introduced and the transition from a signature for natural numbers to a signature for integers is made; the rules in this table extend those of Table 7 and define the canonical term algebra $\mathbb{Z}_{bud}$ that is isomorphic to the canonical term algebra $\mathbb{Z}_{ubd}$ of the specification in Table 3. The DDRS thus defined contains thirty-three (parametric) equations (so, 36+24 eq's in total). We attempt to provide some intuition for equations [b26] and [b27]:

$$(-x):_b i$$

should be equal to $(-x:_b 0) + i$, so $(-x):_b 0 = -(x:_b 0)$, and $(-x):_b 1$ is determined by

$$-(P(x:_b 0)) \overset{[\text{b20}]}{=} -(P(x):_b 1).$$

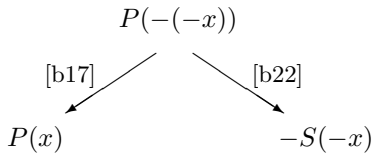
Equations [b24] and [b25] can be explained in a similar way:

$$\begin{array}{lll} S(-(x:_b 0)) & \text{should be equal to} & -(P(x:_b 0)) = -(P(x):_b 1), \\ S(-(x:_b 1)) & \text{should be equal to} & -(P(x:_b 1)) = -(x:_b 0). \end{array}$$

[1] Note that there is no equation [b14.0] that is, $1 = S(0)$, because 1 is a normal form in binary view.

The abbreviation P^j in equations $[\text{b30}.i.j]_{i,j=0}^{1}$ and $[\text{b31}.i.j]_{i,j=0}^{1}$ stands for j applications of the predecessor function P. Normal forms for $\mathbb{Z}_{bud}$ are 0, 1, all applications of $\_:_b 1$ and $\_:_b 1$ to a nonzero normal form, and all minus instances $-t$ for each such normal form t that differs from 0.

We note that the equations in Tables 7 and 8 are semantic consequences of the axioms for commutative rings (equations $(1) - (8)$ in Table 4). The rewriting system defined by this DDRS is proven strongly terminating in [12]. However, its non-confluence is also proven in [12], using the following rewrite steps:



(13)

In Appendix A we prove that this rewriting system for $\mathbb{Z}_{bud}$ is ground-confluent, and thus ground-complete.

[b16]	$-0 = 0$
[b17]	$-(-x) = x$
[b18]	$P(0) = -1$
[b19]	$P(1) = 0$
[b20]	$P(x:_b 0) = P(x):_b 1$
[b21]	$P(x:_b 1) = x:_b 0$
[b22]	$P(-x) = -S(x)$
[b23]	$S(-1) = 0$
[b24]	$S(-(x:_b 0)) = -(P(x):_b 1)$
[b25]	$S(-(x:_b 1)) = -(x:_b 0)$
[b26]	$(-x):_b 0 = -(x:_b 0)$
[b27]	$(-x):_b 1 = -(P(x):_b 1)$
[b28]	$x + (-1) = P(x)$
[b29]	$(-1) + x = P(x)$
$[\text{b30}.i.j]_{i,j=0}^{1}$	$(x:_b i) + (-(y:_b j)) = P^j((x + (-y)):_b i)$
$[\text{b31}.i.j]_{i,j=0}^{1}$	$(-(y:_b j)) + (x:_b i) = P^j((x + (-y)):_b i)$
[b32]	$(-x) + (-y) = -(x + y)$
[b33]	$x \cdot (-y) = -(x \cdot y)$

Table 8: Combined with Table 7, a DDRS for $\mathbb{Z}_{bud}$ that specifies integer numbers in binary view

$[\mathrm{d1}.i]_{i=0}^{9}$	$0:_d i = i$	[d11]	$x \cdot 0 = 0$
$[\mathrm{d2}.i]_{i=0}^{8}$	$S(i) = i'$	$[\mathrm{d12}.i]_{i=0}^{8}$	$x \cdot i' = (x \cdot i) + x$
[d3]	$S(9) = 1:_d 0$	$[\mathrm{d13}.i]_{i=0}^{9}$	$x \cdot (y:_d i) = ((x \cdot y):_d 0) + (x \cdot i)$
$[\mathrm{d4}.i]_{i=0}^{8}$	$S(x:_d i) = x:_d i'$	$[\mathrm{d14}.i]_{i=0}^{1}$	$x:_b i = (x + x) + i$
[d5]	$S(x:_d 9) = S(x):_d 0$		
[d6]	$x + 0 = x$		
[d7]	$0 + x = x$		
$[\mathrm{d8}.i]_{i=1}^{9}$	$x + i = S^i(x)$		
$[\mathrm{d9}.i]_{i=1}^{9}$	$i + x = S^i(x)$		
$[\mathrm{d10}.i.j]_{i,j=0}^{9}$	$(x:_d i) + (y:_d j) = S^j((x + y):_d i)$		

Table 9: A DDRS for $\mathbb{N}_{dub}$, natural numbers in decimal view

2.3 Decimal view

Table 9 defines a DDRS for a decimal view of natural numbers, consisting of fourteen (parametric) equations (172 eq's in total). This DDRS defines the datatype $\mathbb{N}_{dub}$ that is isomorphic to the canonical term algebra $\mathbb{N}_{ubd}$ of the specification in Table 2. Leading zeros except for the zero itself are removed by $[\mathrm{d1}.i]_{i=0}^{9}$, and successor terms are rewritten according to $[\mathrm{d2}.i]_{i=0}^{8} - [\mathrm{d5}]$. Rewriting from binary notation is part of this DDRS, and the last equation scheme $[\mathrm{d14}.i]_{i=0}^{1}$ serves that purpose. The rewriting system defined by the DDRS in Table 9 is proven complete in [12].

Before we extend this DDRS to the integers, we define a variant of successor notation for digits that we call "10 minus" subtraction for decimal digits, and we write

$$i^\star$$

for the "10 minus" decimal digit of i. In Table 10 we display all identities for $i^\star$, and we shall use these in order to cope with terms of the form $(-x):_d i$ for $i = 1, ..., 9$.

$1^\star = 9$	$4^\star = 6$	$7^\star = 3$
$2^\star = 8$	$5^\star = 5$	$8^\star = 2$
$3^\star = 7$	$6^\star = 4$	$9^\star = 1$

Table 10: "10 minus" subtraction for decimal digits

[d15]	$-0 = 0$	$[\text{d27}.i]_{i=1}^{9}$	$x + (-i) = P^i(x)$
[d16]	$-(-x) = x$	$[\text{d28}.i]_{i=1}^{9}$	$(-i) + x = P^i(x)$
[d17]	$P(0) = -1$	$[\text{d29}.i.j]_{i,j=0}^{9}$	$(x :_d i) + (-(y :_d j)) = P^j((x + (-y)) :_d i)$
$[\text{d18}.i]_{i=0}^{8}$	$P(i') = i$	$[\text{d30}.i.j]_{i,j=0}^{9}$	$(-(y :_d j)) + (x :_d i) = P^j((x + (-y)) :_d i)$
[d19]	$P(x :_d 0) = P(x) :_d 9$	[d31]	$(-x) + (-y) = -(x + y)$
$[\text{d20}.i]_{i=0}^{8}$	$P(x :_d i') = x :_d i$	[d32]	$x \cdot (-y) = -(x \cdot y)$
[d21]	$P(-x) = -S(x)$		
$[\text{d22}.i]_{i=0}^{8}$	$S(-i') = -i$		
[d23]	$S(-(x :_d 0)) = -(P(x) :_d 9)$		
$[\text{d24}.i]_{i=0}^{8}$	$S(-(x :_d i')) = -(x :_d i)$		
[d25]	$(-x) :_d 0 = -(x :_d 0)$		
$[\text{d26}.i]_{i=0}^{9}$	$(-x) :_d i = -(P(x) :_d i^\star)$		

Table 11: Combined with Table 9 (and using $i^\star$ from Table 10), a DDRS for $\mathbb{Z}_{dub}$ that specifies integers in decimal view

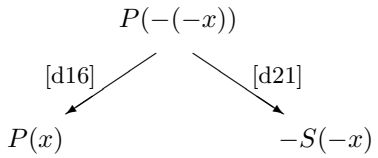
In Table 11 minus and predecessor are added and the transition to a signature for integers is made; the equations in this table extend those of Table 9. The DDRS thus defined is named $\mathbb{Z}_{dub}$ and is isomorphic to the canonical term algebra $\mathbb{Z}_{ubd}$ of the specification in Table 3; it contains thirty-two (parametric) equations (so, $172 + 273$ eq's in total).

The (twenty-one) equations captured by $[\text{d23}] - [\text{d26}.i]_{i=0}^{9}$ can be explained in a similar fashion as was done in the previous section for [b24] − [b27]: for example,

$$(-5) :_d 3$$

should be equal to $-(5 :_d 0) + 3 = -(4 :_d 7)$, and this follows immediately from the appropriate equation in $[\text{d26}.i]_{i=0}^{9}$.

The equations of the DDRS specified by Tables 9 and 11 are semantic consequences of the equations for commutative rings (equations (1) − (8) in Table 4). In [12], the rewriting system defined by this DDRS is proven strongly terminating, and non-confluent by essentially the same counter-example as (13):



We leave it as an open question whether this particular rewriting system is ground-confluent, and thus ground-complete.

3 Alternative DDRSes for integers with digit tree constructors

Having defined DDRSes that employ (postfix) digit append functions in Section 2, we now consider the more general *digit tree constructor* functions. For the binary view, this approach is followed by Bouma and Walters in [7]; for a view based on any radix (number base), this approach is further continued in Walters [14] and Walters and Zantema [15], where the constructor is called *juxtaposition* because it goes with the absence of a function symbol in order to be close to ordinary decimal and binary notation.

We extend the signature $\Sigma_{\mathbb{Z}}$ defined in Section 1.2 with the following three functions (infix):

$$\hat{u}\ ,\ \hat{b}\ ,\ \hat{d}\ : \mathbb{Z} \times \mathbb{Z} \to \mathbb{Z},$$

called "unary digit tree constructor function", "binary digit tree constructor function", and "decimal digit tree constructor function", and to be used for unary, binary notation and decimal notation, respectively. The latter two constructors serve to represent logarithmic notation and satisfy the semantic equations $[\![x\ \hat{b}\ y]\!] = 2 \cdot [\![x]\!] + [\![y]\!]$ and $[\![x\ \hat{d}\ y]\!] = 10 \cdot [\![x]\!] + [\![y]\!]$.

For integer numbers in decimal view or binary view, normal forms are the relevant digits, all applications of the respective constructor with left argument a nonzero normal form and right argument a digit, and all minus instances $-t$ for each such nonzero normal form t, these satisfy $[\![-(t)]\!] = -([\![t]\!])$. E.g.,

$$(9\ \hat{d}\ 7)\ \hat{d}\ 5 \quad \text{and} \quad ((1\ \hat{b}\ 0)\ \hat{b}\ 0)\ \hat{b}\ 1$$

represent the decimal number 975 and the binary number 1001, respectively, and the normal form that represents the additional inverse of the latter is $-(((1\ \hat{b}\ 0)\ \hat{b}\ 0)\ \hat{b}\ 1)$. A minor complication with decimal and binary digit tree constructors is that we now have to consider rewritings such as

$$2\ \hat{d}\ (1\ \hat{d}\ 5) = (2+1)\ \hat{d}\ 5 = 3\ \hat{d}\ 5 \quad (= 35),$$

which perhaps are somewhat non-intuitive. For integers in unary view, thus with unary digit tree constructor, this complication is absent (see Section 3.1).

We keep the presentation of the resulting DDRSes (those defining the binary and decimal view are based on [14, 15]) minimal in the sense that equations for conversion from the one view to the other are left out. Of course, it is easy to define such equations. Also, equations for conversion to and from the datatypes defined in Section 2 are omitted, although such equations are not difficult to define.

3.1 Unary view with digit tree constructor

For naturals in this particular unary view, normal forms are 0 and expressions $t\ \hat{u}\ 0$ with t a normal form (thus, with association of $\hat{u}$ to the left). Of course, the phenomenon of "removing leading zeros" does not exist in this particular unary view (as in the datatype $\mathbb{N}_{u'}$ defined in Table 5). The resulting datatype $\mathbb{N}_{ut}$ is defined in Table 12.

In the unary view, $\hat{u}$ is an associative operator, as is clear from rule [ut1] (in contrast to digit tree constructors for the binary and decimal case). Moreover, the commutative variants $t\ \hat{u}\ r$ and $r\ \hat{u}\ t$ rewrite to the same normal form. The latter property also follows from the following

[ut1]	$x \mathbin{\hat{}_u} (y \mathbin{\hat{}_u} z) = (x \mathbin{\hat{}_u} y) \mathbin{\hat{}_u} z$	[ut4]	$x \cdot 0 = 0$
[ut2]	$x + 0 = x$	[ut5]	$x \cdot (y \mathbin{\hat{}_u} 0) = (x \cdot y) + x$
[ut3]	$x + (y \mathbin{\hat{}_u} 0) = (x + y) \mathbin{\hat{}_u} 0$		

Table 12: A DDRS for $\mathbb{N}_{ut}$, natural numbers in unary view with unary digit tree constructor

semantics for closed terms:

$$\begin{aligned} [\![0]\!] &= 0, \\ [\![x \mathbin{\hat{}_u} y]\!] &= [\![x]\!] + [\![y]\!] + 1, \\ [\![x + y]\!] &= [\![x]\!] + [\![y]\!], \\ [\![x \cdot y]\!] &= [\![x]\!] \cdot [\![y]\!]. \end{aligned}$$

Observe that

$$x + (y \mathbin{\hat{}_u} z) = (x + y) \mathbin{\hat{}_u} z \quad \text{and} \quad x \cdot (y \mathbin{\hat{}_u} z) = (x \cdot (y + z)) + x$$

are valid equations in $\mathbb{N}_{ut}$.

The extension to integer numbers can be done in a similar fashion as in the previous section, thus obtaining normal forms of the form $-(t)$ with t a nonzero normal form in $\mathbb{N}_{ut}$. However, also terms of the form $x \mathbin{\hat{}_u} (-y)$ and variations thereof have to be considered. We define this extension in Table 13 below and call the resulting datatype $\mathbb{Z}_{ut}$.

Adding the interpretation rule $[\![-x]\!] = -[\![x]\!]$ and exploiting the commutativity of $\mathbin{\hat{}_u}$ in $[\![x \mathbin{\hat{}_u} y]\!]$, it can be easily checked that [ut6] – [ut18] (as equations) are sound. Moreover, the equation

$$(-x) \mathbin{\hat{}_u} y = y \mathbin{\hat{}_u} (-x)$$

also holds in $\mathbb{Z}_{ut}$.

[ut6]	$-0 = 0$	[ut13]	$0 + x = x$
[ut7]	$-(-x) = x$	[ut14]	$(x \mathbin{\hat{}_u} 0) + (-(y \mathbin{\hat{}_u} 0)) = x + (-y)$
[ut8]	$0 \mathbin{\hat{}_u} (-(x \mathbin{\hat{}_u} 0)) = -x$	[ut15]	$(-(y \mathbin{\hat{}_u} 0)) + (x \mathbin{\hat{}_u} 0) = x + (-y)$
[ut9]	$(x \mathbin{\hat{}_u} 0) \mathbin{\hat{}_u} (-(y \mathbin{\hat{}_u} 0)) = x \mathbin{\hat{}_u} (-y)$	[ut16]	$(-x) + (-y) = -(x + y)$
[ut10]	$(-(x \mathbin{\hat{}_u} 0)) \mathbin{\hat{}_u} 0 = -x$	[ut17]	$x \cdot (-y) = -(x \cdot y)$
[ut11]	$(-(y \mathbin{\hat{}_u} 0)) \mathbin{\hat{}_u} (x \mathbin{\hat{}_u} 0) = x \mathbin{\hat{}_u} (-y)$	[ut18]	$(-x) \cdot y = -(x \cdot y)$
[ut12]	$(-(x \mathbin{\hat{}_u} 0)) \mathbin{\hat{}_u} (-(y \mathbin{\hat{}_u} 0)) = -((x + y) \mathbin{\hat{}_u} 0)$		

Table 13: Combined with Table 12, a DDRS for $\mathbb{Z}_{ut}$ that specifies integer numbers in unary view with unary digit tree constructor

[bt1]	$0 \mathbin{\hat{b}} x = x$	[bt8]	$x \cdot 0 = 0$
[bt2]	$x \mathbin{\hat{b}} (y \mathbin{\hat{b}} z) = (x + y) \mathbin{\hat{b}} z$	[bt9]	$x \cdot 1 = x$
[bt3]	$0 + x = x$	[bt10]	$x \cdot (y \mathbin{\hat{b}} z) = (x \cdot y) \mathbin{\hat{b}} (x \cdot z)$
[bt4]	$1 + 0 = 1$		
[bt5]	$1 + 1 = 1 \mathbin{\hat{b}} 0$		
[bt6]	$1 + (x \mathbin{\hat{b}} y) = x \mathbin{\hat{b}} (1 + y)$		
[bt7]	$(x \mathbin{\hat{b}} y) + z = x \mathbin{\hat{b}} (y + z)$		

Table 14: A DDRS for $\mathbb{N}_{bt}$, natural numbers in binary view with binary digit tree constructor

3.2 Binary view with digit tree constructor

For naturals in binary view with the binary digit tree constructor, the associated datatype $\mathbb{N}_{bt}$ is defined in Table 14. According to [15] (with a reference to [7]), the rewriting system defined by [bt1] – [bt7] is strongly terminating and ground-confluent, and thus ground-complete.

In [15] a rewriting system for integer arithmetic is provided with next to juxtaposition and (unary) minus also addition, subtraction and multiplication, and proven ground-confluent and terminating with respect to any radix (number base). In Table 15 we present a variant of this rewriting system without subtraction for the binary digit tree constructor, and define the datatype $\mathbb{Z}_{bi}$.

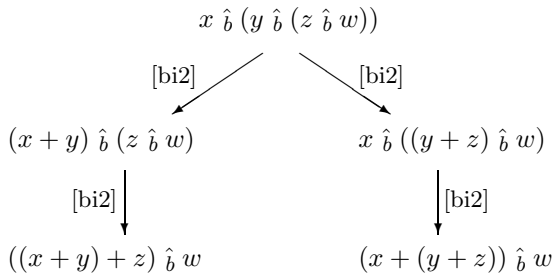
[bi1]	$0 \mathbin{\hat{b}} x = x$	[bi13]	$-0 = 0$
[bi2]	$x \mathbin{\hat{b}} (y \mathbin{\hat{b}} z) = (x + y) \mathbin{\hat{b}} z$	[bi14]	$-(-x) = x$
[bi3]	$0 + x = x$	[bi15]	$1 \mathbin{\hat{b}} (-1) = 1$
[bi4]	$x + 0 = x$	[bi16]	$(x \mathbin{\hat{b}} 0) \mathbin{\hat{b}} (-1) = (x \mathbin{\hat{b}} (-1)) \mathbin{\hat{b}} 1$
[bi5]	$1 + 1 = 1 \mathbin{\hat{b}} 0$	[bi17]	$(x \mathbin{\hat{b}} 1) \mathbin{\hat{b}} (-1) = (x \mathbin{\hat{b}} 0) \mathbin{\hat{b}} 1$
[bi6]	$x + (y \mathbin{\hat{b}} z) = y \mathbin{\hat{b}} (x + z)$	[bi18]	$x \mathbin{\hat{b}} (-(y \mathbin{\hat{b}} z)) = -((y + (-x)) \mathbin{\hat{b}} z)$
[bi7]	$(x \mathbin{\hat{b}} y) + z = x \mathbin{\hat{b}} (y + z)$	[bi19]	$(-x) \mathbin{\hat{b}} y = -(x \mathbin{\hat{b}} (-y))$
[bi8]	$x \cdot 0 = 0$	[bi20]	$1 + (-1) = 0$
[bi9]	$0 \cdot x = 0$	[bi21]	$(-1) + 1 = 0$
[bi10]	$1 \cdot 1 = 1$	[bi22]	$x + (-(y \mathbin{\hat{b}} z)) = -(y \mathbin{\hat{b}} (z + (-x)))$
[bi11]	$x \cdot (y \mathbin{\hat{b}} z) = (x \cdot y) \mathbin{\hat{b}} (x \cdot z)$	[bi23]	$(-(x \mathbin{\hat{b}} y)) + z = -(x \mathbin{\hat{b}} (y + (-z)))$
[bi12]	$(x \mathbin{\hat{b}} y) \cdot z = (x \cdot z) \mathbin{\hat{b}} (y \cdot z)$	[bi24]	$x \cdot (-y) = -(x \cdot y)$
		[bi25]	$(-x) \cdot y = -(x \cdot y)$

Table 15: A DDRS for $\mathbb{Z}_{bi}$, integer numbers in binary view with binary digit tree constructor

[dt1]	$0 \mathbin{\hat{d}} x = x$	[dt7]	$x + 0 = x$
[dt2]	$x \mathbin{\hat{d}} (y \mathbin{\hat{d}} z) = (x + y) \mathbin{\hat{d}} z$	$[\text{dt8}.i]_{i=0}^{8}$	$x + i' = S(x) + i$
$[\text{dt3}.i]_{i=0}^{8}$	$S(i) = i'$	$[\text{dt9}.i]_{i=0}^{9}$	$x + (y \mathbin{\hat{d}} i) = (y \mathbin{\hat{d}} x) + i$
[dt4]	$S(9) = 1 \mathbin{\hat{d}} 0$	[dt10]	$x \cdot 0 = 0$
$[\text{dt5}.i]_{i=0}^{8}$	$S(x \mathbin{\hat{d}} i) = x \mathbin{\hat{d}} i'$	$[\text{dt11}.i]_{i=0}^{8}$	$x \cdot i' = x + (x \cdot i)$
[dt6]	$S(x \mathbin{\hat{d}} 9) = S(x) \mathbin{\hat{d}} 0$	$[\text{dt12}.i]_{i=0}^{9}$	$x \cdot (y \mathbin{\hat{d}} i) = ((x \cdot y) \mathbin{\hat{d}} 0) + (x \cdot i)$

Table 16: A DDRS for $\mathbb{N}_{dt}$, natural numbers with decimal digit tree constructor in decimal view (using i' from Table 1)

In [12] it is proven that the associated term rewriting system is strongly terminating. Confluence is disproven in [12] by the following counter-example:



(14)

3.3 Decimal view with digit tree constructor

For naturals in decimal view with the decimal digit tree constructor, we make use of successor terms, in order to avoid (non-parametric) equations such as

$$1 + 1 = 2, ..., \quad 9 + 8 = 1 \mathbin{\hat{d}} 7, \quad 9 + 9 = 1 \mathbin{\hat{d}} 8,$$
$$1 \cdot 1 = 1, ..., \quad 8 \cdot 9 = 7 \mathbin{\hat{d}} 2, \quad 9 \cdot 9 = 8 \mathbin{\hat{d}} 1.$$

The associated datatype $\mathbb{N}_{dt}$ is defined in Table 16. Note that equations of the form

$$i' + x = i + S(x)$$

instead of (or next to) $[\text{dt8}.i]_{i=0}^{8}$ would destroy termination: $2+1 \mapsto 1+S(1) \mapsto S(1)+1 \mapsto 2+1$. Moreover, the interplay between digit tree constructor, successor and normal form notation makes it by equations $[\text{dt9}.i]_{i=0}^{9}$ possible not to incorporate the equation $0 + x = x$ in this particular, relatively simple DDRS. According to [12], the associated rewriting system is strongly terminating, but not confluent (cf. counter-example (14)).

The extension to integers is given by the equations in Table 17 that define the datatype $\mathbb{Z}_{dt}$. In contrast to the approaches in [14, 15] with juxtaposition, we now make use of both successor

[dt1]	$0 \mathbin{\hat{d}} x = x$	[dt13]	$-0 = 0$
[dt2]	$x \mathbin{\hat{d}} (y \mathbin{\hat{d}} z) = (x + y) \mathbin{\hat{d}} z$	[dt14]	$-(-x) = x$
$[\text{dt3}.i]_{i=0}^{8}$	$S(i) = i'$	[dt15]	$P(0) = -1$
[dt4]	$S(9) = 1 \mathbin{\hat{d}} 0$	$[\text{dt16}.i]_{i=0}^{8}$	$P(i') = i$
$[\text{dt5}.i]_{i=0}^{8}$	$S(x \mathbin{\hat{d}} i) = x \mathbin{\hat{d}} i'$	[dt17]	$P(x \mathbin{\hat{d}} 0) = P(x) \mathbin{\hat{d}} 9$
[dt6]	$S(x \mathbin{\hat{d}} 9) = S(x) \mathbin{\hat{d}} 0$	$[\text{dt18}.i]_{i=0}^{8}$	$P(x \mathbin{\hat{d}} i') = x \mathbin{\hat{d}} i$
[dt7]	$x + 0 = x$	[dt19]	$P(-x) = -S(x)$
$[\text{dt8}.i]_{i=0}^{8}$	$x + i' = S(x) + i$	$[\text{dt20}.i]_{i=0}^{8}$	$S(-i') = -i$
$[\text{dt9}.i]_{i=0}^{9}$	$x + (y \mathbin{\hat{d}} i) = (y \mathbin{\hat{d}} x) + i$	[dt21]	$S(-(x \mathbin{\hat{d}} 0)) = -(P(x) \mathbin{\hat{d}} 9)$
[dt10]	$x \cdot 0 = 0$	$[\text{dt22}.i]_{i=0}^{8}$	$S(-(x \mathbin{\hat{d}} i')) = -(x \mathbin{\hat{d}} i)$
$[\text{dt11}.i]_{i=0}^{8}$	$x \cdot i' = x + (x \cdot i)$	[dt23]	$(-x) \mathbin{\hat{d}} y = -(x \mathbin{\hat{d}} (-y))$
$[\text{dt12}.i]_{i=0}^{9}$	$x \cdot (y \mathbin{\hat{d}} i) = ((x \cdot y) \mathbin{\hat{d}} 0) + (x \cdot i)$	$[\text{dt24}i.j]_{i,j=1}^{9}$	$i \mathbin{\hat{d}} (-j) = P(i) \mathbin{\hat{d}} j^{\star}$
		$[\text{dt25}.i]_{i=1}^{9}$	$(x \mathbin{\hat{d}} y) \mathbin{\hat{d}} (-i) = P(x \mathbin{\hat{d}} y) \mathbin{\hat{d}} i^{\star}$
		[dt26]	$x \mathbin{\hat{d}} (-(y \mathbin{\hat{d}} z)) = -((y + (-x)) \mathbin{\hat{d}} z)$
		$[\text{dt27}.i]_{i=1}^{9}$	$x + (-i) = P^i(x)$
		[dt28]	$x + (-(y \mathbin{\hat{d}} z)) = -(y \mathbin{\hat{d}} (z + (-x)))$
		$[\text{dt29}.i]_{i=1}^{9}$	$(-i) + x = P^i(x)$
		[dt30]	$(-(x \mathbin{\hat{d}} y)) + z = -(x \mathbin{\hat{d}} (y + (-z)))$
		[dt31]	$x \cdot (-y) = -(x \cdot y)$
		[dt32]	$(-x) \cdot y = -(x \cdot y)$

Table 17: A DDRS for $\mathbb{Z}_{dt}$, integer numbers with decimal digit tree constructor in decimal view (using i' from Table 1 and $i^{\star}$ from Table 10)

terms and predecessor terms, and the DDRS presented here is composed from rewrite rules for successor and predecessor, rewrite rules defined in [14, 15], and combinations thereof. Note that we can still do without the equation $0 + x = x$:

$0 + (-0) = 0$	by [dt13] and [dt7],
$0 + (-i) = P^i(0) = ... = -i$	by $[\text{dt27}.i]_{i=1}^{9}$ (and some more equations),
$0 + (-(t_1 \mathbin{\hat{d}} t_2)) = -(t_1 \mathbin{\hat{d}} t_2)$	by equations [dt28], [dt13], and [dt7].

As stated in [12], it is an open question whether this term rewriting system for $\mathbb{Z}_{dt}$ is strongly terminating and/or ground-confluent.

4 Concluding remarks

This paper is about the design (by means of trial and error) of *datatype defining rewrite systems* (DDRSes) rather than about the precise analysis of the various rewriting systems per se. What matters in addition to readability and conciseness of each DDRS is at this stage a reasonable confidence that each of these rewriting systems is strongly terminating and ground-confluent (and thus ground-complete), and that the intended normal forms are reached by means of rewriting.

When specifying a datatype of integers as an extension of the naturals, the unary view leads to satisfactory results, but with high inefficiency. For the binary view and the decimal view based on the unary append functions and discussed in Section 2, corresponding extensions are provided, but the resulting rewriting systems are at first sight significantly less concise and comprehensible. Recently, strong termination has been automatically proven by Kluiving and van Woerkom [12] with the use of the AProVE tool [11]. Some further remarks:

1. The three DDRSes (datatype defining rewrite systems) for integers given in Section 2 each produce an extension datatype for a datatype for the natural numbers. An initial algebra specification of the datatype of integers is obtained from any of the DDRSes given in [1] by

   - taking the reduct to the signature involving unary, binary, and decimal notation only,
   - removing rewrite rules involving operators for hexadecimal notation,
   - expanding the signature with a unary additive inverse and a unary predecessor function,
   - adding rewrite rules (in equational form) that allow for the unique normalization of closed terms involving the minus sign,

   while making sure that these rewrite rules (viewed as equations) are semantic consequences of the equations for commutative rings.

2. Syntax for hexadecimal notation has been omitted because that usually plays no role when dealing with integers. It is an elementary exercise to incorporate hexadecimal notation.

3. The DDRSes for the binary view and the decimal view are hardly intelligible unless one knows that the objective is to construct a commutative ring. A decimal normal form is defined as either a digit, or an application of a decimal append function $\_ :_d i$ to a nonzero normal form (for all digits i). This implies the absence of (superfluous) leading zeros, and the (ground) normal forms thus obtained correspond bijectively to the non-negative integers (that is, $\mathbb{N}$). Incorporating all minus instances $-(t)$ for each nonzero normal form t yields the class of normal forms. The "semantics" of these normal forms in the language of commutative rings is very simple:

$$\begin{aligned} [\![0]\!] &= 0, \\ [\![i']\!] &= [\![i]\!] + 1 \quad \text{for all digits } i \text{ (and } i' \text{ defined as in Table 1)}, \\ [\![x :_d i]\!] &= (\underline{10} \cdot [\![x]\!]) + [\![i]\!] \quad \text{for all digits } i \text{ and } \underline{10} = [\![9]\!] + 1, \\ [\![-(x)]\!] &= -([\![x]\!]). \end{aligned}$$

   A binary normal form has similar semantics: $[\![x :_b i]\!] = (\underline{2} \cdot [\![x]\!]) + [\![i]\!]$ for digits $0, 1$, and $\underline{2} = 1 + 1$.

4. Understanding the concept of a commutative ring can be expected only from a person who has already acquired an understanding of the structure of integers and who accepts the concept of generalization of a structure to a class of structures sharing some but not all of its properties.

   In other words, the understanding that a DDRS for the integers is provided in the binary view and in the decimal view can only be communicated to an audience under the assumption that a reliable mental picture of the integers already exists in the minds of members of the audience. This mental picture, however, can in principle be communicated by taking notice of the DDRS for the unary view first.

   This conceptual (near) circularity may be nevertheless be considered a significant weakness of the approach of defining (and even introducing) the integers as an extension of naturals by means of rewriting.

Although full confluence of the DDRSes in Section 2 has been disproven by KLUIVING and VAN WOERKOM [12] (with the use of CSI [17]), we prove that the DDRS that defines the binary view of the integers (Tables 7 and 8) is ground-complete by proving its ground-confluence in Appendix A.

In Section 3 we discussed some alternatives for the above-mentioned DDRSes based on papers of BOUMA and WALTERS [7], WALTERS [14], and WALTERS and ZANTEMA [15] in which digit tree constructors are employed. In this case, a digit is a normal form, and so is an application of the digit tree constructor that adheres to association to the left and with the removal of (superfluous) leading zeros. Thus, $n \mathbin{\hat{d}} i$ is a normal form if n is a nonzero normal form and i a digit. Incorporating all minus instances $-(t)$ for each nonzero normal form t yields the class of normal forms for integers. With the tool AProVE [11], KLUIVING and VAN WOERKOM [12] proved strong termination of the DDRSes that employ the binary tree constructor for $\mathbb{N}$ (Table 14) and for $\mathbb{Z}$ (Table 15), and for the DDRS that uses the decimal tree constructor for $\mathbb{N}$ (Table 16), but had to leave this question open for the case of the decimal tree constructor for $\mathbb{Z}$ (Table 17). For all these DDRSes, full confluence was automatically disproven in [12] (using the CSI tool [17]).

Of course, a decimal notation as 689 is so common that one usually does not question whether it represents $(6 :_d 8) :_d 9$ or $(6 \mathbin{\hat{d}} 8) \mathbin{\hat{d}} 9$ or some other formally defined notation. Nevertheless, as we have seen, different algorithmic approaches to for example addition may apply, although one would preferably not hamper an (initial) arithmetical method with notation such as $x \mathbin{\hat{d}} (y \mathbin{\hat{d}} z)$ and rewrite rules such as $x \mathbin{\hat{d}} (y \mathbin{\hat{d}} z) = (x + y) \mathbin{\hat{d}} z$ and those for $+$, and for this reason we have a preference for the DDRSes defined in Section 2.

In [14], WALTERS presents a TRS (term rewriting system) based on juxtaposition for integer arithmetic with addition and subtraction that is ground-complete; this system is proven ground-confluent and terminating with respect to any radix. In [15], WALTERS and ZANTEMA extend this system with multiplication and prove ground-completeness, using *semantic labelling* for their termination proof, and judge this TRS to have good efficiency and readability (in comparison with some alternatives discussed in that paper). Furthermore, the authors also discuss a TRS that is based on successor and predecessor notation, and in which minus is not used: negative numbers are represented by normal forms $P(0), P(P(0))$, and so on. This TRS is comparable to the DDRS in Table 3 that defines $\mathbb{Z}_{ubd}$ and is proven confluent and terminating, and judged to have poor readability and (too) high complexity. Finally, a complete TRS for natural numbers with addition and multiplication based on digit append is also provided in [15].

We briefly mention another, comparable approach to integer arithmetic that is also based on some form of digit append constructors for representing integer numbers. In [9], CONTEJEAN, MARCHÉ and RABEHASAINA introduce integer arithmetic based on *balanced ternary numbers*, that

[r1]	$-0 = 0$	[r8]	$0 + x = x$
[r2]	$-(-x) = x$	[r9]	$(-1) + 1 = 0$
		[r10]	$(-(x + 1)) + 1 = -x$
[r3]	$x + (y + z) = (x + y) + z$	[r11]	$(-x) + (-y) = -(x + y)$
[r4]	$x + 0 = x$		
[r5]	$1 + (-1) = 0$	[r12]	$x \cdot 0 = 0$
[r6]	$(x + 1) + (-1) = x$	[r13]	$x \cdot 1 = x$
[r7]	$x + (-(y + 1)) = (x + (-y)) + (-1)$	[r14]	$x \cdot (-y) = (-x) \cdot y$
		[r15]	$x \cdot (y + z) = (x \cdot y) + (x \cdot z)$

Table 18: A DDRS for the ring of integers

is, numbers that can be represented by a digit append function $:_t$ with digits -1,0,1 and semantics $[\![i]\!] = i$ and $[\![x :_t i]\!] = \underline{3} \cdot [\![x]\!] + i$ (see, e.g., KNUTH [13]) and provide a TRS that is confluent and terminating modulo associativity and commutativity of addition and multiplication.

Based on either a DDRS for the natural numbers or a DDRS for the integers one may develop a DDRS for rational numbers in various ways. It is plausible to consider the meadow of rational numbers of [6] or the non-involutive meadow of rational numbers (see [2]) or the common meadow of rational numbers (see [3]) as abstract algebraic structures for rationals in which unary, binary, and decimal notation are to be incorporated in ways possibly based on the specifications presented above. Furthermore, one does well to consider the work discussed in [9] on a term rewriting system for rational numbers, in which arithmetic for rational numbers is specified (this is the main result in [9], for which the above-mentioned work on integer arithmetic is a preliminary): the authors specify rational numbers by means of a TRS that is complete modulo associativity and commutativity of addition and multiplication, taking advantage of Stein's algorithm for computing gcd's of non-negative integers without any division[2] (see, e.g., [13]).

A survey of equational algebraic specifications for abstract datatypes is provided in [16]. In [5] one finds the general result that computable abstract datatypes can be specified by means of specifications which are confluent and strongly terminating term rewriting systems. Some general results on algebraic specifications can be found in [8, 4, 10]. More recent applications of equational specifications can be found in [6].

We conclude the paper with the introduction of a very simple DDRS that specifies the integers in the signature $\Sigma_r = \{0, 1, -(\_), +, \cdot\}$ of rings. This DDRS is defined in Table 18. Observe that the negative variant of equation [r7], that is,

$$(-x) + (y + 1) = ((-x) + y) + 1$$

is an instance of equation [r3]. Also, observe that the equations in Table 18 are semantic consequences of the axioms for commutative rings (equations $(1) - (8)$ in Table 4). In [12], KLUIVING and VAN WOERKOM report that the term rewriting system defined by this DDRS is strongly terminating,[3] and below we prove that it is also ground-confluent, and thus ground-complete.

[2] Apart from halving even numbers, which is easy in binary notation, but can otherwise be specified with a shift operation.

[3] Alternatively, the following weight function $|t|$ on closed terms can be used to prove strong termination: $|0| = |1| = 2$, $|x + y| = |x| + 2|y|$, $|-x| = 1 + 3/2|x|$, and $|x \cdot y| = |x| \cdot |y|^2$.

Define the set NF of closed terms over Σ_r as follows:

$$\begin{aligned} NF &= \{0\} \cup NF^+ \cup NF^-, \\ NF^+ &= \{1\} \cup \{t+1 \mid t \in NF^+\}, \\ NF^- &= \{-t \mid t \in NF^+\}. \end{aligned}$$

It immediately follows that if $t \in NF$, then t is a normal form (no rewrite step applies). Furthermore, two distinct elements in NF have distinct values in $\mathbb{Z}$. In order to prove ground-confluence of the associated TRS we have to show that for each closed term over Σ_r, either $t \in NF$ or t has a rewrite step, so that each normal form is in NF.

We prove this by structural induction on t. The base cases $t \in \{0, 1\}$ are trivial. For the induction step we have to consider the following cases:

- Case $t = -r$. Assume that $r \in NF$ and apply case distinction on r:
  - if $r = 0$, then $t \to 0$ by equation [r1],
  - if $r \in NF^+$, then $t \in NF$,
  - if $r \in NF^-$, then t has a rewrite step by equation [r2].
- Case $t = u + r$. Assume that $u, r \in NF$ and apply case distinction on r:
  - if $r = 0$, then $t \to u$ by equation [r4],
  - if $r = 1$, then apply case distinction on u:
    - if $u = 0$, then $t \to 1$ by equation [r8],
    - if $u \in NF^+$, then $t \in NF$,
    - if $u = -1$, then $t \to 0$ by equation [r9],
    - if $u = -(u' + 1)$, then t has a rewrite step by equation [r10],
  - if $r = r' + 1$, then $t \to (u + r') + 1$ by equation [r3],
  - if $r = -1$ then $t = u + (-1)$ and apply case distinction on u:
    - if $u = 0$, then t has a rewrite step by equation [r8],
    - if $u = 1$, then t has a rewrite step by equation [r5],
    - if $u = u' + 1$, then t has a rewrite step by equation [r6],
    - if $u \in NF^-$, then t has a rewrite step by equation [r11],
  - if $r = -(r' + 1)$, then $t \to (u + (-r')) + (-1)$ by equation [r7].
- Case $t = u \cdot r$. Assume that $u, r \in NF$, then t has a rewrite step according to one of the equations [r12] – [r15].

This concludes our proof.

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Appendix A

We prove that the term rewriting system defined by the DDRS for $\mathbb{Z}_{bud}$ in Tables 7 and 8 is ground-complete. This rewriting system is proven strongly terminating in [12], so it remains to be proven that it is ground-confluent and we adopt the approach used in the proof on page 19.

Define the set N of closed terms over $\Sigma_{\mathbb{Z}}$ as follows:

$$\begin{aligned} N &= \{0\} \cup N^+ \cup N^-, \\ N^+ &= \{1\} \cup \{w :_b 0, w :_b 1 \mid w \in N^+\}, \\ N^- &= \{-t \mid t \in N^+\}. \end{aligned}$$

It immediately follows that if $t \in N$, then t is a normal form (no rewrite rule applies), and that two distinct elements in N have distinct values in $\mathbb{Z}$. Also, as stated in Section 2.2, the equations in Tables 7 and 8 are semantic consequences of the axioms for commutative rings (equations $(1) - (8)$ in Table 4). In order to prove ground-confluence of this rewriting system we have to show that for each closed term t in $\mathbb{Z}_{bud}$, either $t \in N$ or t has a rewrite step, so that each normal form is in N. We prove this by structural induction on t.

The base cases are simple: if $t \in \{0, 1\}$, then $t \in N$, and if $t \in \{3, 4, 5, 6, 7, 8, 9\}$, then t has a rewrite step according to one of $[\text{b14}.i]_{i=1}^{8}$. For the induction step we have to consider eight cases:

1. Case $t = -r$. Assume that $r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to 0$ by equation [b16],
   - if $r = 1$, then $t \in N$,
   - if $r = -1$, then $t \to 1$ by equation [b17],
   - if $r = v :_b i$, then $t \in N$,
   - if $r = -(v :_b i)$, then $t \to v :_b i$ by equation [b17].
2. Case $t = S(r)$. Assume that $r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to 1$ by equation [b2],
   - if $r = 1$, then $t \to 1 :_b 0$ by equation [b3],
   - if $r = -1$, then $t \to 0$ by equation [b23],
   - if $r = v :_b 0$, then $t \to v :_b 1$ by equation [b4],
   - if $r = v :_b 1$, then $t \to S(v) :_b 0$ by equation [b5],
   - if $r = -(v :_b 0)$, then $t \to -(P(v) :_b 1)$ by equation [b24],

- if $r = -(v :_b 1)$, then $t \to -(v :_b 0)$ by equation [b25].

3. Case $t = P(r)$. Assume that $r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to -1$ by equation [b18],
   - if $r = 1$, then $t \to 0$ by equation [b19],
   - if $r = -1$, then $t \to -S(1)$ by equation [b22],
   - if $r = v :_b 0$, then $t \to P(v) :_b 1$ by equation [b20],
   - if $r = v :_b 1$, then $t \to v :_b 0$ by equation [b21],
   - if $r = -(v :_b i)$, then $t \to -S(v :_b i)$ by equation [b22].
4. Case $t = r :_b 0$. Assume that $r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to 0$ by the first equation of $[\text{b1}.i]_{i=0}^{1}$,
   - if $r = 1$ or $r = v :_b i$, then $t \in N$,
   - if $r = -1$ or $r = -(v :_b i)$, then t has a rewrite step by equation [b26].
5. Case $t = r :_b 1$. Assume that $r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to j$ by the second equation of $[\text{b1}.i]_{i=0}^{1}$,
   - if $r = 1$ or $r = v :_b i$, then $t \in N$,
   - if $r = -1$ or $r = -(v :_b i)$, then t has a rewrite step by equation [b27].
6. Case $t = r \mathbin{\hat{d}} i$ with i a digit (thus, $i \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$). Now t has a rewrite step according to one of the equations of $[\text{b15}.i]_{i=0}^{9}$.
7. Case $t = u + r$. Assume that $u, r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to u$ by equation [b6],
   - if $r = 1$, then $t \to S(u)$ by equation [b8],
   - if $r = -1$, then $t \to P(u)$ by equation [b28].
   - if $r = v :_b i$, apply a case distinction on u:
     - if $u = 0$, then $t \to r$ by equation [b7],
     - if $u = 1$, then $t \to S(r)$ by equation [b9],
     - if $u = -1$, then $t \to P(r)$ by equation [b29],
     - if $u = w :_b j$, then t has a rewrite step according to one of $[\text{b10}.i.j]_{i,j=0}^{1}$,
     - if $u = -(w :_b j)$, then t has a rewrite step according to one of $[\text{b30}.i.j]_{i,j=0}^{1}$,
   - if $r = -(v :_b i)$, apply a case distinction on u:
     - if $u = 0$, then $t \to r$ by equation [b7],
     - if $u = 1$, then $t \to S(r)$ by equation [b9],
     - if $u = -1$, then $t \to P(r)$ by equation [b29],
     - if $u = w :_b j$, then t has a rewrite step according to one of $[\text{b31}.i.j]_{i,j=0}^{1}$,
     - if $u = -(w :_b j)$, then t has a rewrite step by equation [b32].
8. Case $t = u \cdot r$. Assume that $u, r \in N$ and apply case distinction on r:
   - if $r = 0$, then $t \to 0$ by equation [b11],
   - if $r = 1$, then $t \to u$ by equation [b12],
   - if $r = -1$ or $r = -(v :_b i)$, then t has a rewrite step by equation [b33].
   - if $r = v :_b i$, then t has a rewrite step according to one of $[\text{b13}.i]_{i=0}^{1}$.

This concludes our proof.

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